

# Homework 1

DUE: Wednesday, 8. Sept. 2021

**Exercise 1.** By construction, categories that arise from posets have the following property: between any given pair of objects, there is at most one morphism. Does every category with this property arise from a poset? What additional properties are needed to characterize those categories that arise from posets?

**Exercise 2.** A 2-category  $\mathbb{C}$  consists of

- a set  $\text{ob}(\mathbb{C})$  of objects,
- for every pair  $(x, y)$  of objects, a category  $\mathbb{C}(x, y)$  of morphisms from  $x$  to  $y$ ,
- for every object  $x$  an object  $\text{id}_x \in \mathbb{C}(x, x)$  called the identity morphism,
- for every triple  $x, y, z$ , a functor

$$\mu : \mathbb{C}(x, y) \times \mathbb{C}(y, z) \rightarrow \mathbb{C}(x, z)$$

called a composition law,

subject to the conditions

1. for every objects  $x, y \in \mathbb{C}$ , the functors

$$\mu(-, \text{id}_y) : \mathbb{C}(x, y) \rightarrow \mathbb{C}(x, y)$$

and

$$\mu(\text{id}_x, -) : \mathbb{C}(x, y) \rightarrow \mathbb{C}(x, y)$$

are the identity functors on  $\mathbb{C}(x, y)$ ,

2. for every 4-tuple  $(x, y, z, w)$  of objects of  $\mathbb{C}$ , the diagram

$$\begin{array}{ccc} \mathbb{C}(x, y) \times \mathbb{C}(y, z) \times \mathbb{C}(z, w) & \xrightarrow{\mu \times \text{id}} & \mathbb{C}(x, z) \times \mathbb{C}(z, w) \\ \downarrow \text{id} \times \mu & & \downarrow \mu \\ \mathbb{C}(x, y) \times \mathbb{C}(y, w) & \xrightarrow{\mu} & \mathbb{C}(x, w) \end{array}$$

commutes.

Show that the set of (small) categories forms the objects of a 2-category  $\text{Cat}$  where, for small categories  $\mathbb{C}, \mathbb{D}$ , we define  $\text{Cat}(\mathbb{C}, \mathbb{D}) = \text{Fun}(\mathbb{C}, \mathbb{D})$ .

**Exercise 3.** Suppose  $F : \mathbb{C} \longrightarrow \mathbb{D}$  and  $G : \mathbb{D} \longrightarrow \mathbb{C}$  are adjoint functors with unit  $\epsilon : \text{id}_{\mathbb{C}} \Rightarrow G \circ F$  and counit  $\eta : F \circ G \Rightarrow \text{id}_{\mathbb{D}}$ .

1. Show that  $\epsilon$  is a natural isomorphism if and only if  $F$  is fully faithful.
2. Show that  $\eta$  is a natural isomorphism if and only if  $G$  is fully faithful.

**Exercise 4.** For a category  $\mathbf{C}$ , we will use the notation  $\mathbf{Set}_{\mathbf{C}} := \mathbf{Fun}(\mathbf{C}^{\text{op}}, \mathbf{Set})$ .

1. Show that the Yoneda embedding

$$\mathcal{Y} : \mathbf{C} \longrightarrow \mathbf{Set}_{\mathbf{C}}$$

defined in class is a functor.

2. Use the Yoneda lemma to show that  $\mathcal{Y}$  is fully faithful.
3. Is  $\mathcal{Y}$  essentially surjective?